

Nonadditive quantum error-correcting code

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We report the first nonadditive quantum error-correcting code, namely, a $((9, 12, 3))$ code which is a 12-dimensional subspace within a 9-qubit Hilbert space, that outperforms the optimal stabilizer code of the same length by encoding more levels while correcting arbitrary single-qubit errors.

The quantum error-correcting code (QECC) [1, 2, 3, 4] provides an active way of protecting our quantum data from decohering. Almost all the QECCs constructed so far are stabilizer codes [5, 6, 7], codes that have the structure of an eigenspace of an Abelian group generated by multilocal Pauli operators. Codes without such a structure are called nonadditive codes. The first nonadditive code [8, 9] that outperforms the stabilizer codes is the $((5, 6, 2))$ code, a 5-qubit code encoding 6 levels capable of correcting single-qubit *erasure*, i.e., a code of distance 2. Recently a family of distance 2 nonadditive codes with a higher encoding rate has been constructed [10]. Though some nonadditive error-correcting codes had been constructed [11, 12], the question of whether the nonadditive error-correcting codes with a distance larger than 2 can encode more levels than the corresponding stabilizer codes remains open.

In this Letter we report the first nonadditive code of distance 3 that beats the corresponding stabilizer code: a nonadditive $((9, 12, 3))$ code that is a 12-dimensional subspace in a 9-qubit Hilbert space against arbitrary single-qubit errors. In comparison, the best stabilizer code $[[9, 3, 3]]$ of the same length can encode only 3 logical qubits, i.e., an 8-dimensional subspace [7].

Our new code is most conveniently formulated in terms of graph states [13, 14]. Let $G = (V, \Gamma)$ be an undirected simple graph with $|V| = n$ vertices and Γ , called as the *adjacency matrix* of the graph, is an $n \times n$ symmetric matrix with vanishing diagonal entries and $\Gamma_{ab} = 1$ if vertices a, b are connected and $\Gamma_{ab} = 0$ otherwise. Consider a system of n qubits labeled by V and denote by $\mathcal{X}_a, \mathcal{Y}_a$, and \mathcal{Z}_a three Pauli operators acting on qubit $a \in V$. The *graph state* associated with graph G reads

$$|G\rangle = \prod_{ab=1} \mathcal{U}_{ab} |+\rangle_x^V = \frac{1}{\sqrt{2^n}} \sum_{\vec{\mu}=0}^1 (-1)^{\frac{1}{2}\vec{\mu} \cdot \Gamma \cdot \vec{\mu}} |\vec{\mu}\rangle_z, \quad (1)$$

where $|\vec{\mu}\rangle_z$ is the common eigenstates of $\{\mathcal{Z}_a\}_{a \in V}$ with $(-1)^{\mu_a}$ as eigenvalues, $|+\rangle_x^V$ denotes the simultaneous +1 eigenstate of $\{\mathcal{X}_a\}_{a \in V}$, and $\mathcal{U}_{ab} = (1 + \mathcal{Z}_a + \mathcal{Z}_b - \mathcal{Z}_a \mathcal{Z}_b)/2$ is the controlled-phase operation between qubit a and b . The graph state is also the unique simultaneous +1 eigenstate of n vertex stabilizers $\mathcal{G}_a = \mathcal{X}_a \mathcal{Z}_{N_a}$ with $a \in V$ where N_a is the neighborhood of a and we denote by $\mathcal{Z}_U = \prod_{a \in U} \mathcal{Z}_a$ for a subset of vertices $U \subseteq V$.

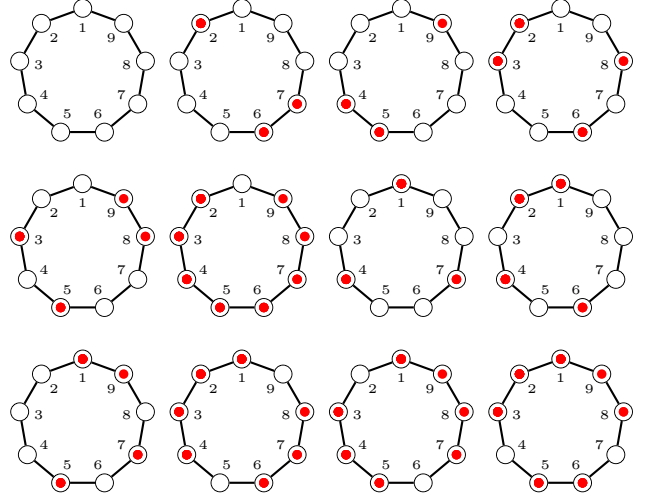


FIG. 1: (Color online) Twelve graph-state bases on the loop graph L_9 for the $((9, 12, 3))$ code \mathbb{D} . Each graph represents a graph state that is the unique common eigenstate of the vertex stabilizers $\{\mathcal{G}_a\}$ with eigenvalue +1 if a is uncolored and -1 if the vertex is red-colored.

We consider in what follows the loop graph L_9 on 9 vertices which are labeled by integers from 1 to 9. Its adjacency matrix has nonvanishing entries $\Gamma_{aa\pm} = 1$ ($1 \leq a \leq 9$) only where $a_{\pm} = a \pm 1$ with identifications $9_+ = 1$ and $1_- = 9$. The corresponding graph state is denoted as $|L_9\rangle$. We claim that the 12-dimensional subspace \mathbb{D} spanned by the states $\{\mathcal{Z}_{V_i} |L_9\rangle\}_{i=1}^{12}$ where

$$\begin{aligned} V_1 &= \emptyset, V_2 = \{2, 6, 7\}, V_3 = \{4, 5, 9\}, V_4 = \{2, 3, 6, 8\} \\ V_5 &= \{3, 5, 8, 9\}, V_6 = \{2, 3, 4, 5, 6, 7, 8, 9\} \\ V_7 &= \{1, 4, 7\}, V_8 = \{1, 2, 4, 6\}, V_9 = \{1, 5, 7, 9\} \\ V_{10} &= \{1, 2, 3, 4, 6, 7, 8\}, V_{11} = \{1, 3, 4, 5, 7, 8, 9\} \\ V_{12} &= \{1, 2, 3, 5, 6, 8, 9\}, \end{aligned} \quad (2)$$

as shown in Fig.1, is a $((9, 12, 3))$ code. Obviously these 12 states are mutually orthogonal since V_i 's are distinct and $\langle G | \mathcal{Z}_V | G \rangle = \delta_{V, \emptyset}$ holds true for any graph state. To prove that the code is of distance 3, i.e., capable of correcting single-qubit errors, we have only to demonstrate that each one of 3×9 single-qubit errors and 9×36 two-qubit errors \mathcal{E} will bring \mathbb{D} into its orthogonal complement [4, 13], i.e.,

$$\langle L_9 | \mathcal{Z}_{V_i} \mathcal{E} \mathcal{Z}_{V_j} | L_9 \rangle = 0, \quad (1 \leq i, j \leq 12). \quad (3)$$

Since all the bases of \mathbb{D} given above are the common eigenstates of the vertex stabilizers $\{\mathcal{G}_a = \mathcal{Z}_{a-}\mathcal{X}_a\mathcal{Z}_{a+}\}_{a=1}^9$, a bit flip error \mathcal{X}_a on these bases is equivalent to a phase flip error \mathcal{Z}_{N_a} on qubits in its neighborhood, e.g., $N_a = \{a_+, a_-\}$ in L_9 , upto an unimportant phase factor. And a \mathcal{Y}_a error can be equivalently replaced by a phase flip error $\mathcal{Z}_a\mathcal{Z}_{N_a}$ on qubits a, a_+ , and a_- . As a result every single-qubit error is equivalent to one of the following phase flip errors $\{\mathcal{Z}_a, \mathcal{Z}_{N_a}, \mathcal{Z}_a\mathcal{Z}_{N_a}\}$ for $1 \leq a \leq 9$ and every two-qubit error is equivalent to one of the following phase flip errors

$$\begin{array}{lll} \mathcal{Z}_a\mathcal{Z}_b, & \mathcal{Z}_{N_a}\mathcal{Z}_b, & \mathcal{Z}_{N_a}\mathcal{Z}_a\mathcal{Z}_b, \\ \mathcal{Z}_a\mathcal{Z}_{N_b}, & \mathcal{Z}_{N_a}\mathcal{Z}_{N_b}, & \mathcal{Z}_a\mathcal{Z}_{N_a}\mathcal{Z}_{N_b}, \\ \mathcal{Z}_a\mathcal{Z}_b\mathcal{Z}_{N_b}, & \mathcal{Z}_{N_a}\mathcal{Z}_{N_b}\mathcal{Z}_b, & \mathcal{Z}_a\mathcal{Z}_b\mathcal{Z}_{N_a}\mathcal{Z}_{N_b}, \end{array} \quad (4)$$

with $1 \leq a, b \leq 9$. To summarize, for a loop graph, every single-qubit or two-qubit error is equivalent to one of following 6 patterns of phase flip errors

$$\begin{array}{ll} \text{I:} & \mathcal{Z}_a, \\ \text{II:} & \mathcal{Z}_a\mathcal{Z}_b, \\ \text{III:} & \mathcal{Z}_{a-}\mathcal{Z}_b\mathcal{Z}_{a+}, \mathcal{Z}_{a\pm}\mathcal{Z}_a\mathcal{Z}_{a\pm 3}, \\ \text{IV:} & \mathcal{Z}_{a-}\mathcal{Z}_{a+}\mathcal{Z}_{b-}\mathcal{Z}_{b+}, \mathcal{Z}_{a-}\mathcal{Z}_a\mathcal{Z}_{a+}\mathcal{Z}_b, \\ & \mathcal{Z}_{a-}\mathcal{Z}_{a-2}\mathcal{Z}_{a+}\mathcal{Z}_{a+2}, \\ \text{V:} & \mathcal{Z}_{a-}\mathcal{Z}_a\mathcal{Z}_{a+}\mathcal{Z}_{b-}\mathcal{Z}_{b+}, \\ \text{VI:} & \mathcal{Z}_{a-}\mathcal{Z}_a\mathcal{Z}_{a+}\mathcal{Z}_{b-}\mathcal{Z}_b\mathcal{Z}_{b+}, \end{array} \quad (5)$$

where a, b are suitably chosen so that error patterns I, II, III, IV, V, VI are phase flip errors on 1 qubit to 6 qubits respectively. It is clear that phase flip errors on more than 6 qubits cannot be caused by any single-qubit or two-qubit error.

As an immediate consequence, Eq.(3) is equivalent to saying that *none* of the transition operators $\mathcal{Z}_{V_i}\mathcal{Z}_{V_j}$ ($1 \leq i < j \leq 12$) between each pair of bases of \mathbb{D} belongs to any one of 6 error patterns listed in Eq.(5). Because $\mathcal{Z}_{V_k}\mathcal{Z}_{V_7} = \mathcal{Z}_{V_{k+6}}$ it is enough to examine the following 31 different transition operators

$$\begin{array}{llllll} \mathcal{Z}_{147}, \mathcal{Z}_{126}, \mathcal{Z}_{1246}, \mathcal{Z}_{2368}, \mathcal{Z}_{12569}, & \mathcal{Z}_{1234678}, \mathcal{Z}_{12345689}, \\ \mathcal{Z}_{159}, \mathcal{Z}_{1348}, \mathcal{Z}_{2569}, \mathcal{Z}_{23678}, & \mathcal{Z}_{1235689}, \mathcal{Z}_{12356789}, \\ \mathcal{Z}_{267}, \mathcal{Z}_{1378}, \mathcal{Z}_{3589}, \mathcal{Z}_{34589}, & \mathcal{Z}_{1245679}, \mathcal{Z}_{23456789}, \\ \mathcal{Z}_{348}, \mathcal{Z}_{1579}, & \mathcal{Z}_{123468}, \mathcal{Z}_{1345789}, \\ \mathcal{Z}_{378}, \mathcal{Z}_{2467}, & \mathcal{Z}_{135789}, \mathcal{Z}_{2345689}, \\ \mathcal{Z}_{459}, \mathcal{Z}_{4579}, & \mathcal{Z}_{245679}, \mathcal{Z}_{2356789}. \end{array}$$

obtained from \mathcal{Z}_{V_7} and $\{\mathcal{Z}_{V_i}\mathcal{Z}_{V_j}, \mathcal{Z}_{V_7}\mathcal{Z}_{V_i}\mathcal{Z}_{V_j} | 1 \leq i < j \leq 6\}$. It is easy to check that phase flip errors on 5 or more qubits in the above table do not belong to any one of the error patterns in Eq.(5). Because of the symmetry of the loop graph L_9 , one needs only to check that \mathcal{Z}_{126} , \mathcal{Z}_{147} , \mathcal{Z}_{1246} , and \mathcal{Z}_{2368} do not belong to any one of the error patterns in Eq.(5), which are easy tasks to perform. In this way we have demonstrated that \mathbb{D} is a $((9, 12, 3))$ code, which is obviously nonadditive.

As to the projector of the code \mathbb{D} , we notice that there are 3 local stabilizers of the code \mathbb{D} , namely, \mathcal{G}_{38} , \mathcal{G}_{62} , and \mathcal{G}_{95} , where we have denoted $\mathcal{G}_U = \prod_{v \in U} \mathcal{G}_v$ for a subset

of vertices U . By denoting

$$\begin{aligned} \mathcal{A} = & \mathcal{G}_{14} \left(1 - \mathcal{G}_{36} + \mathcal{G}_{39} - \mathcal{G}_{69} + 2\mathcal{G}_{369} + 2\mathcal{G}_9 \right) \\ & + \mathcal{G}_{17} \left(1 - \mathcal{G}_{39} + \mathcal{G}_{36} - \mathcal{G}_{69} + 2\mathcal{G}_{369} + 2\mathcal{G}_6 \right), \end{aligned} \quad (6)$$

we can write down the projector of the code \mathbb{D} as

$$\mathcal{P} = \frac{1}{2^{10}} (1 + \mathcal{G}_{38})(1 + \mathcal{G}_{62})(1 + \mathcal{G}_{95})\mathcal{A}(\mathcal{A} + 8), \quad (7)$$

from which the weight enumerator [15, 16] of the code \mathbb{D} can be readily obtained

$$2^9 \times 12 \times \left(\frac{3u^9}{128} + \frac{u^5v^4}{64} + \frac{u^3v^6}{4} + \frac{u^2v^7}{2} + \frac{27uv^8}{128} \right). \quad (8)$$

Here the coefficients of $u^{9-d}v^d$ is given by $\sum \text{Tr}^2(\mathcal{P}\mathcal{E}_d)$ with the summation taken over all (Hermitian) errors acting nontrivially on d qubits.

To conclude, we have provided the first evidence that nonadditive error-correcting codes can perform better than the stabilizer codes. Since the bases of our code are all graph states, they can be easily be prepared from a product state by using controlled phase operation and local unitary operations as shown in Eq.(1).

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